A Context-free Grammar for the Ramanujan-Shor Polynomials

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Dedicated to Joseph P. S. Kung

Abstract

Ramanujan defined the polynomials $\psi_k(r,x)$ in his study of power series inversion. Berndt, Evans and Wilson obtained a recurrence relation for $\psi_k(r,x)$. In a different context, Shor introduced the polynomials Q(i, j, k) related to improper edges of a rooted tree, leading to a refinement of Cayley's formula. He also proved a recurrence relation and raised the question of finding a combinatorial proof. Zeng realized that the polynomials of Ramanujan coincide with the polynomials of Shor, and that the recurrence relation of Shor coincides with the recurrence relation of Berndt, Evans and Wilson. So we call these polynomials the Ramanujan-Shor polynomials, and call the recurrence relation the Berndt-Evans-Wilson-Shor recursion. A combinatorial proof of this recursion was obtained by Chen and Guo, and a simpler proof was recently given by Guo. From another perspective, Dumont and Ramamonjisoa found a context-free grammar G to generate the number of rooted trees on n vertices with k improper edges. Based on the grammar G, we find a grammar H for the Ramanujan-Shor polynomials. This leads to a formal calculus for the Ramanujan-Shor polynomials. In particular, we obtain a grammatical derivation of the Berndt-Evans-Wilson-Shor recursion. We also provide a grammatical approach to the Abel identities and a grammatical explanation of the Lacasse identity.

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1 Introduction

For integers $1 \le k \le r + 1$, Ramanujan [14] defined the polynomials $\psi_k(r, x)$ by the following relation:

$$\sum_{k=0}^{\infty} \frac{(x+k)^{r+k} e^{-u(x+k)} u^k}{k!} = \sum_{k=1}^{r+1} \frac{\psi_k(r,x)}{(1-u)^{r+k}},\tag{1.1}$$

and derived the recurrence relation:

$$\psi_k(r,x) = (x-1)\psi_k(r,x-1) + \psi_{k-1}(r+1,x) - \psi_{k-1}(r+1,x-1), \tag{1.2}$$

where $\psi_1(0,x) = 1$, $\psi_0(r,x) = 0$ and $\psi_k(r,x) = 0$ for k > r + 1. Berndt, Evans and Wilson [1] obtained another recurrence relation

$$\psi_k(r,n) = (n-r-k+1)\psi_k(r-1,n) + (r+k-2)\psi_{k-1}(r-1,n).$$
 (1.3)

By setting u = 0 in (1.1), Ramanujan deduced the identity

$$\sum_{k=1}^{r+1} \psi_k(r, x) = x^r. \tag{1.4}$$

Zeng [20] observed that the polynomials $\psi_k(r,x)$ coincide with the polynomials introduced by Shor [17]. Let

$$Q_{n,k}(x) = \psi_{k+1}(n-1, x+n). \tag{1.5}$$

Then (1.4) can be rewritten as

$$\sum_{k=0}^{n-1} Q_{n,k}(x) = (x+n)^{n-1}, \tag{1.6}$$

and (1.3) can be recast as

$$Q_{n,k}(x) = (x-k+1)Q_{n-1,k}(x+1) + (n+k-2)Q_{n-1,k-1}(x+1).$$
(1.7)

Shor [17] defined the polynomials $Q_{n,k}(x)$ in the notation of Q(i,j,k). To be more specific, the numbers Q(i,j,k) are defined by the recurrence relation

$$Q(i,j,k) = (k+i-1)Q(i-1,j,k) + (i+j-2)Q(i-1,j-1,k),$$
(1.8)

where Q(1,0,k) = 1, Q(i,-1,k) = 0 for $i \ge 1$ and Q(1,j,k) = 0 for $j \ge 1$.

Let [n] denote the set $\{1, 2, ..., n\}$. Shor showed that, for $k \geq 1$, kQ(i - k, j, k) equals the number of forests on [i] rooted at $\{1, 2, ..., k\}$ with j improper edges. For a rooted tree T and a vertex of v of T, we use T_v to denote the subtree of T rooted at v, namely, the subtree of T consisting of the vertex v together with all the descendants of v. An edge (u, v) of T, with v being the child of u, is called an improper edge if there exists a vertex in T_v that is smaller than v. This leads to a combinatorial proof of the relation

$$\sum_{i=0}^{i-1} Q(i,j,k) = (i+k)^{i-1}, \tag{1.9}$$

which can be considered as a refinement of Cayley's formula. As noted by Shor [17], Q(i, j, k) is a polynomial in k for $i \ge 1$ and $0 \le j \le i - 1$, and (1.9) holds when k is replaced by a variable x. Indeed, the polynomials Q(i, j, x) can be expressed as $Q_{i,j}(x)$, and so we call $Q_{n,k}(x)$ the Ramanujan-Shor polynomials. Notice that in the notation of $Q_{n,k}(x)$, (1.9) turns to be exactly the formula (1.6).

Shor proved that in addition to the recurrence relation (1.8), the polynomials Q(i, j, k) also satisfy the following relation: For $i \ge 1$, $0 \le j \le i - 1$ and $k \ge 1$,

$$Q(i,j,k) = (k-j+1)Q(i-1,j,k+1) + (i+j-2)Q(i-1,j-1,k+1).$$
 (1.10)

He asked the question of finding a combinatorial interpretation of the above recurrence relation. Notice that (1.10) is precisely the recursion (1.7) proved by Berndt, Evans and Wilson. So we call it the Berndt-Evans-Wilson-Shor recursion. It is worth mentioning that the recursion (1.8) of Shor can be deduced from the recursion (1.2) of Ramanujan and the Berndt-Evans-Wilson-Shor recursion (1.7).

Zeng [20] found the following interpretations of the polynomials $Q_{n,k}(x)$ in terms of the number of improper edges of trees on [n+1] with root 1:

$$Q_{n,k}(x) = \sum_{T \in F_{n+1,k}} x^{\deg_T(1)-1}, \tag{1.11}$$

where $F_{n,k}$ denotes the set of trees on [n] with k improper edges and with root 1, and $\deg_T(1)$ denotes the degree of the vertex 1 in T. Zeng also showed that the polynomials $Q_{n,k}(x)$ can be interpreted by the number of improper edges of rooted trees (not necessarily rooted at 1) on [n], namely,

$$Q_{n,k}(x) = \sum_{T \in R_{n,k}} (x+1)^{\deg_T(1)}, \tag{1.12}$$

where $R_{n,k}$ denotes the set of rooted trees on [n] with k improper edges. In answer to a question of Zeng [20], Chen and Guo [4] found a bijection showing that (1.11) and (1.12) are equivalent.

We should also note that (1.11) is equivalent to the interpretation of $Q_{n,k}(x)$ given by Shor when x is a positive integer. As noted by Shor [17], for a positive integer r, $rQ_{n,k}(r)$ equals the number of forests on [n+r] rooted at $\{1,2,\ldots,r\}$ with a total number of k improper edges. Let F be such a forest counted by $rQ_{n,k}(r)$. Let T_i be the tree in F rooted at i, where $1 \leq i \leq r$. For each T_i , removing the root i and coloring the subtrees of T_i with color i, we get a forest on $\{r+1,r+2,\ldots,r+n\}$ with each tree colored by one of colors $1,2,\ldots,r$. After relabeling, this leads to a forest on [n] with each tree associated with one of the colors $1,2,\ldots,r$. Let $U_{n,k}$ denote the set of forests of rooted trees on [n] with k improper edges. For a forest F in $U_{n,k}$, let tree (F) denote the number of trees in F. By the above argument, one sees that

$$rQ_{n,k}(r) = \sum_{F \in U_{n,k}} r^{\operatorname{tree}(F)},$$

which is equivalent to (1.11) since a forest F in $U_{n,k}$ gives rise to a rooted tree T in $F_{n+1,k}$ by adding a new root 0.

Subtracting (1.7) from (1.8), the Berndt-Evans-Wilson-Shor recursion (1.7) takes the form

$$Q_{n,k}(1+x) = Q_{n,k}(x) + (n+k-1)Q_{n-1,k}(1+x),$$
(1.13)

where $n \geq 1, 0 \leq k \leq n-1$, $Q_{1,0}(x) = 1$ and $Q_{n,k}(x) = 0$ if $k \geq n$ or k < 0. Chen and Guo [4] gave a combinatorial proof of (1.13) in answer to the question of Shor. More precisely, let $F_{n+1,k}[\deg(2) > 0]$ denote the set of rooted trees in $F_{n+1,k}$ in which the vertex 2 is not a leaf, and let $F_{n+1,k}[\deg(n+1) > 0]$ denotes the set of rooted trees in $F_{n+1,k}$ in which the vertex n+1 is not a leaf. A bijection between $F_{n+1,k}[\deg(2) > 0]$ and $F_{n+1,k}[\deg(n+1) > 0]$ was constructed in [4]. A simpler bijection was given by Guo [9].

Based on Shor's recursive procedure to construct rooted trees, Dumont and Ramamonjisoa [7] found a context-free grammar to enumerate rooted trees with a given

number of improper edges. They defined a grammar G by the following substitution rules:

$$G: A \to A^3S, S \to AS^2.$$

Let D denote the formal derivative with respect to G. Dumont and Ramamonjisoa showed that, for $n \geq 1$,

$$D^{n-1}(AS) = A^n S^n \sum_{k=0}^{n-1} b(n,k) A^k,$$

where b(n, k) denotes the number of rooted trees on [n] with k improper edges. Note that $b(n, k) = Q_{n,k}(0)$.

Based on the Dumont-Ramamonjisoa grammar, we obtain a grammar H to generate the Ramanujan-Shor polynomials $Q_{n,k}(x)$. Let

$$H: a \to axy, x \to xyw, y \to y^3w, w \to yw^2,$$

and let D denote the formal derivative with respect to H. For $n \geq 1$, we obtain the following relation

$$D^{n}(a) = axy^{n}w^{n-1} \sum_{k=0}^{n-1} Q_{n,k}(xw^{-1})y^{k}.$$

With the aid of the grammar H, we are led to a simple derivation of the Berndt-Evans-Wilson-Shor recursion in the form of (1.13).

It turns out that the grammar H can also be used to deal with the Abel identities. In a certain sense, the formal derivative with respect to the grammar H can be viewed as a shift-invariant operator for the Abel identities in the spirit of the umbral calculus, see Rota [16]. As will be seen, the Abel identities can be deduced from the Leibnitz formula with respect to the grammar H.

Riordan [15] defined the sum

$$A_n(x_1, x_2; p, q) = \sum_{k=0}^{n} \binom{n}{k} (x_1 + k)^{k+p} (x_2 + n - k)^{n-k+q},$$

where $n \geq 1$ and the parameters p,q are integers. He found closed formulas of $A_n(x_1,x_2;p,q)$ for some p and q. These identities were called the Abel identities or the Abel-type identities since the case (p,q)=(-1,0) corresponds to the classical Abel identity. We give a grammar H' based on the grammar H and show that the summations $A_n(x_1,x_2;p,q)$ can be evaluated by using the grammar H'. Using this approach, closed forms can be deduced for $A_n(x,y;-1,0)$, $A_n(x_1,x_2,-1,-1)$ and $A_n(x_1,x_2,-2,0)$ and $A_n(x_1,x_2;-2,-2)$. The case for $A_n(x_1,x_2;-2,-2)$ seems to be new.

We conclude this paper with a grammatical explanation of the identity:

$$n^{n+1} = \sum_{k=1}^{n} \sum_{k=0}^{n-j} \binom{n}{j} \binom{n-j}{k} j^{j} k^{k} (n-j-k)^{n-j-k}.$$

We call this identity the Lacasse identity. It was conjectured by Lacasse [10] in the study of the PAC-Bayesian machine learning theory. Since then, several proofs have been found. For example, Sun [18] gave a derivation by using the umbral calculus, Younsi [19] found a proof with aid of the Abel identity, Prodinger [13] provided a justification based on Cauchy's integral formula, Gessel [8] proved the identity by means of the Lagrange inversion formula, and Chen, Peng and Yang [5] obtained a combinatorial interpretation in terms of triply rooted trees.

This paper is organized as follows. In Section 2, we give an overview of the Dumont-Ramamonjisoa grammar and introduce a grammatical labeling of labeled trees. In Section 3, we find a grammar H to generate the Ramanujan-Shor polynomials. Section 4 is devoted to a proof of the Berndt-Evans-Wilson-Shor relation by using the grammar H. In Section 5, we consider the grammatical derivations of Abel identities. We also provide a grammatical explanation of the Lacasse identity.

2 The Dumont-Ramamonjisoa Grammar

In this section, we give an overview of the context-free grammar introduced by Dumont and Ramamonjisoa [7] to generate rooted trees. The approach of using context-free grammars to combinatorial polynomials was introduced in [2]. Further studies can be found in [3,6,7,11,12]. A context-free grammar G over an alphabet A is defined to be a set of production rules. Given a context-free grammar, one may define a formal derivative D as a differential operator on polynomials or Laurent polynomials in A, that is, D is a linear operator satisfying the relation

$$D(uv) = D(u)v + uD(v),$$

and in general the Leibnitz formula

$$D^{n}(uv) = \sum_{k=0}^{n} \binom{n}{k} D^{k}(u) D^{n-k}(v).$$
 (2.1)

Dumont and Ramamonjisoa [7] defined the following grammar

$$G: A \to A^3 S, S \to A S^2.$$
 (2.2)

Let D denote the formal derivative with respect to the grammar G. Notice that D can also be viewed as the operator

$$D = A^3 S \frac{\partial}{\partial A} + A S^2 \frac{\partial}{\partial S}.$$

Dumont and Ramamonjisoa established a connection between the grammar G and the enumeration of rooted trees on [n] with k improper edges. The notion of an improper edge of a rooted tree was introduced by Shor. Let T a rooted tree on [n]. An edge of T is represented by a pair (u, v) of vertices with v being a child of u. We say that an edge (u, v) of T is improper if v < u or there exists a descendant of v that is smaller than u; otherwise, (u, v) is called a proper edge. Recall that b(n, k) denotes the number of rooted trees on [n] with k improper edges. Dumont and Ramamonjisoa obtained the following relation.

Theorem 2.1 For $n \geq 1$,

$$D^{n-1}(AS) = A^n S^n \sum_{k=0}^{n-1} b(n,k) A^k.$$
 (2.3)

For example, for n = 1, 2, 3, we have

$$D^{0}(AS) = AS,$$

$$D^{1}(AS) = D(A)S + AD(S) = A^{2}S^{2}(1+A),$$

$$D^{2}(AS) = D(D(AS)) = A^{3}S^{3}(2+4A+3A^{2}).$$

Dumont and Ramamonjisoa gave a proof of the above theorem by showing that the coefficients of $D^n(AS)$ satisfy the recurrence relation (1.8) of Shor. More precisely, let s(n,k) denote the coefficient of $A^{n+k}S^n$ in $D^{n-1}(AS)$, Dumont and Ramamonjisoa proved that

$$s(n,k) = (n-1)s(n-1,k) + (n+k-2)s(n-1,k-1),$$

which is equivalent to the relation (1.8) for the case x = 0.

Here we present a proof in the language of a grammatical labeling of rooted trees, which was introduced in [3]. Let R_n denote the set of rooted trees on [n] and let F_n denote the set of rooted trees on [n] with root 1. Recall that $F_{n,k}$ is the set of rooted trees in F_n with k improper edges, and $R_{n,k}$ is the set of rooted trees in R_n with k improper edges. Shor [17] provided a construction of a rooted tree in R_n by adding the vertex n into a tree in R_{n-1} . For a better understanding of the construction, let

us consider the following procedure to delete the vertex n from a rooted tree T in R_n to obtain a rooted tree T' in R_{n-1} . For a rooted tree $T \in R_n$ and a vertex u in T, we adopt the notation $\beta_T(u)$, or simply $\beta(u)$, for the minimum vertex among the vertices in the subtree of T rooted at u.

- 1. Case 1: n is a leaf in T. Delete the vertex n.
- 2. Case 2: n is not a leaf. Assume that n has t children b_1, b_2, \ldots, b_t . We may further assume that

$$\beta(b_1) < \beta(b_2) < \cdots < \beta(b_t).$$

Contract the edge (n, b_t) and relabel the resulting vertex by b_t .

Conversely, one can construct a rooted tree T on [n] with k or k+1 improper edges from a rooted tree T' on [n-1] with k improper edges. There are four operations to add the vertex n to T'.

- 1. Adding n to the tree T' as a child of an arbitrary vertex v, we obtain a tree $T \in R_{n,k}$ with n being a leaf.
- 2. Splitting a proper edge (i, j) into (i, n) and (n, j), we obtain a tree $T \in R_{n,k+1}$. In this case, the degree of n equals one.
- 3. Splitting an improper edge (i, j) into (i, n) and (n, j), we also obtain a tree $T \in R_{n,k+1}$. In this case, the degree of n also equals one.
- 4. Choose an improper edge (v, b_j) in T', where v has t children b_1, b_2, \ldots, b_t listed in the increasing order of their β -values. We relabel v by n and make v a child of n. Moreover, assign b_1, \ldots, b_j to be the children of n and assign b_{j+1}, \ldots, b_t to be the children of v. Then we are led to a tree $T \in R_{n,k+1}$. In this case, the degree of n in T is at least two.

As will be seen, the above construction is closely related to the grammar G. To demonstrate this connection, we introduce a grammatical labeling of rooted trees. We may view a rooted tree T on [n] as a rooted tree \hat{T} on $\{0, 1, 2, \ldots, n\}$ with 0 being the root with only one child. Clearly, the edge below the root 0 of \hat{T} is always a proper edge. Moreover, we represent an improper edge by double edges, called the left edge and the right edge. The idea of using double edges to represent an improper edge is due to Dumont and Ramamonjisoa [7]. We label a vertex of \hat{T} except for 0 by S and label an edge of \hat{T} by A. In other words, a proper edge of T is labeled by A and an improper edge of T is labeled by A^2 . The weight of T is defined by the product of the

labels attached to \hat{T} , denoted by w(T). Apparently, for any tree T in $R_{n,k}$, we have $w(T) = A^{n+k}S^n$.

Figure 2.1 illustrates all rooted trees on $\{1,2,3\}$, where the improper edges are represented by double edges, and the vertex 0 is added at the top of each tree in R_3 .

$$\begin{bmatrix} 0 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 2 & 2 & 3 & 3 \\ 2 & 3 & 2 & 3 & 1 & 2 \end{bmatrix}$$

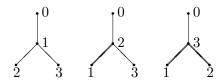


Figure 2.1: Rooted trees in R_3 .

The following relation is a restatement of Theorem 2.1.

Theorem 2.2 For
$$n \ge 1$$
,
$$D^{n-1}(AS) = \sum_{T \in R_n} w(T). \tag{2.4}$$

In view of the above grammatical laleling of rooted trees, it can be seen that the four cases in Shor's construction of a tree T' on [n] from a tree on [n-1] correspond to the substitution rules in G. Instead of giving a detail proof, let us use an example to demonstrate the correspondence.

In Figure 2.2, T is a rooted tree on $\{1, 2, 3, 4\}$. The weight of T is $w(T) = A^6 S^4$. The trees T_1, T_2, T_3 and T_4 are obtained from T in the four cases of Shor's construction.

Case 1: T_1 is obtained from T by adding the vertex 5 as a leaf. Comparing the weights of T_1 and T, it can be seen that this operation corresponds to the substitution rule $S \to AS^2$. Notice that the label S indicates where one can apply this operation.

Case 2: T_2 is obtained from T by splitting the left edge (4,1) into two edges (4,5) and (5,1). This operation corresponds to the substitution rule $A \to A^3S$.

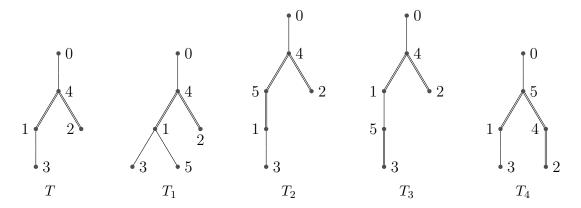


Figure 2.2: An example for the operator D

Case 3: T_3 is obtained from T by splitting the proper edge (1,3) into (1,5) and (5,3). This operation also corresponds to the substitution rule $A \to A^3S$.

Case 4: T_4 is obtained by adding 5 to T via the following procedure: 4 is relabeled by 5, a new vertex 4 is added as a child of 5, the subtree rooted by 1 and the subtree rooted by 2 are assigned as a child of 5 and a child of 4, respectively. It can be seen that this operation also corresponds to the substitution rule $A \to A^3S$.

The above argument is sufficient to lead to a rigorous proof of relation (2.4).

3 A grammar for the R-S polynomials

In this section, we give a grammar H to generate the Ramanujan-Shor polynomials $Q_{n,k}(x)$. Define

$$H: a \to axy, x \to xyw, y \to y^3w, w \to yw^2.$$
 (3.1)

Recall that a rooted tree $T \in F_{n,k}$ is rooted at 1 and has k improper edges. For $T \in F_{n,k}$, we label a proper edge by y, and represent each improper edge of T by double edges, which are both labeled by y. Meanwhile, we label the vertex 1 by a, label each child of the vertex 1 by x and label other vertices by x, so that for $x \in F_{n+1,k}$, the weight of $x \in F_{n+1,k}$ the

$$w(T) = ax^{\deg_T(1)}y^{n+k}w^{n-\deg_T(1)}. (3.2)$$

Figure 3.3 illustrates a rooted tree in $F_{6,2}$ with weight $w(T) = ax^2y^7w^3$.

Let D denote the formal derivative with respect to the grammar H. Recall that F_n is the set of rooted tree on [n] with root 1. The next theorem shows that D can be used to generate the sum of weights of rooted trees in F_n .

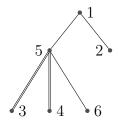


Figure 3.3: A rooted tree $T \in F_{6,2}$

Theorem 3.1 For $n \ge 1$,

$$D^n(a) = \sum_{T \in F_n} w(T). \tag{3.3}$$

To prove the above relation, it is sufficient to observe that the substitution rules in H correspond to the changes of labels in Shor's construction according to the above labeling scheme.

Figure 3.4 gives three rooted trees T_1 , T_2 and T_3 obtained from the tree T in Figure 3.3 by adding the vertex 7 as a leaf as in Case 1 of Shor's construction. Since 7 is child of the root 1, $w(T_1)$ is obtained from w(T) by applying the substitution rule $a \to axy$. Similarly, 7 is a child of 2 in T_2 , and $w(T_2)$ is obtained from w(T) by utilizing the rule $x \to xyw$. Since 7 is a child of 6 in T_3 , $w(T_3)$ is obtained from w(T) by the rule $w \to yw^2$.

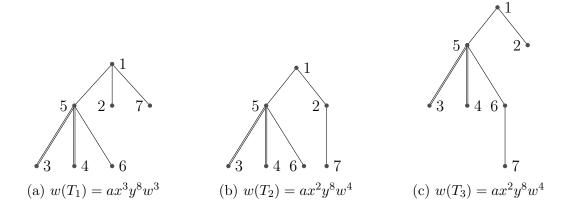


Figure 3.4: The action of D

For Case 2, Case 3 and Case 4 in Shor's construction, the changes of weights of the resulting trees can be characterized by the rule $y \to y^3 w$, just like the rule $A \to A^3 S$ in the Dumont-Ramamonjisoa grammar.

Let $Q_n(x,y)$ denote the generating function of $Q_{n,k}(x)$, that is,

$$Q_n(x,y) = \sum_{k=0}^{n-1} Q_{n,k}(x)y^k = \sum_{T \in F_{n+1}} x^{\deg_T(1)-1} y^{\operatorname{imp}(T)}.$$

For n = 1, 2, 3, we have

$$Q_1(x, y) = 1.$$

 $Q_2(x, y) = y + x + 1.$
 $Q_3(x, y) = 3y^2 + (3x + 4)y + x^2 + 3x + 2.$

We now come to a relationship between the grammar H and the polynomials $Q_n(x,y)$.

Theorem 3.2 For $n \ge 1$,

$$D^{n}(a) = axy^{n}w^{n-1}Q_{n}(xw^{-1}, y). (3.4)$$

For n = 1, 2, 3, we have

$$D(a) = axy = axyQ_0(xw^{-1}, y),$$

$$D^2(a) = ax^2y^2 + axy^2w + axy^3w = axy^2w(xw^{-1} + 1 + y) = axy^2wQ_1(xw^{-1}, y),$$

$$D^3(a) = ax^3y^3 + 3ax^2y^3w + 3ax^2y^4w + 2axy^3w^2 + 4axy^4w^2 + 3axy^5w^2$$

$$= axy^3w^2(x^2w^{-2} + 3xw^{-1} + 2 + (4xw^{-1} + 3)y + 3y^2)$$

$$= axy^3w^2Q_2(xw^{-1}, y).$$

We end this section with a grammatical derivation of the relation (1.8) of Shor, which can be restated below in the notation of $Q_{n,k}(x)$.

Theorem 3.3 For $n \ge 2$ and $1 \le k \le n-2$, we have

$$Q_{n,k}(x) = (x+n-1)Q_{n-1,k}(x) + (n+k-2)Q_{n-1,k-1}(x).$$
(3.5)

Proof. For $n \geq 1$, by the definition of $Q_n(x,y)$, (3.4) can be written as

$$D^{n}(a) = (xw^{-1})ay^{n}w^{n} \sum_{k=0}^{n-1} y^{k}Q_{n,k}(xw^{-1}).$$
(3.6)

For $n \geq 2$, substituting n by n - 1, (3.6) takes the form

$$D^{n-1}(a) = (xw^{-1})ay^{n-1}w^{n-1}\sum_{k=0}^{n-2} y^k Q_{n-1,k}(xw^{-1}).$$
(3.7)

Since

$$D(xw^{-1}) = xyw \cdot w^{-1} - x \cdot w^{-2}yw^2 = 0, (3.8)$$

that is, xw^{-1} is a constant with respect to D, we find that

$$D\left(y^{k}Q_{n-1,k}(xw^{-1})\right) = Q_{n-1,k}(xw^{-1})D(y^{k}) = ky^{k+2}wQ_{n-1,k}(xw^{-1}).$$

Meanwhile,

$$D(ay^{n-1}w^{n-1}) = axy^n w^{n-1} + (n-1)ay^{n+1}w^n + (n-1)ay^n w^n.$$

Therefore, applying the operator D to both sides of (3.7) yields

$$D^{n}(a) = (xw^{-1})ay^{n}w^{n} \left\{ \sum_{k=0}^{n-2} ky^{k+1}Q_{n-1,k}(xw^{-1}) + (xw^{-1} + (n-1)y + (n-1)) \sum_{k=0}^{n-2} y^{k}Q_{n-1,k}(xw^{-1}) \right\}.$$
(3.9)

For $n \ge 2$ and $0 \le k \le n-2$, comparing the coefficients of $axy^{n+k}w^n$ on the right-hand sides of (3.6) and (3.9), we deduce that

$$Q_{n,k}(xw^{-1}) = (xw^{-1} + n - 1)Q_{n-1,k}(xw^{-1}) + (n+k-2)Q_{n-1,k-1}(xw^{-1}).$$
 (3.10)

Setting w = 1 completes the proof.

4 The Berndt-Evans-Wilson-Shor Recursion

The section is devoted to a grammatical derivation of the Berndt-Evans-Wilson-Shor recurrence relation (1.7). To this end, we establish a grammatical expression for $Q_n(r+x,y)$, where r is a nonnegative integer.

Theorem 4.1 For $n \ge 1$ and $r \ge 0$,

$$D^{n}(ax^{r}) = ax^{r}y^{n}w^{n}(r + xw^{-1})Q_{n}(r + xw^{-1}, y).$$
(4.1)

To prove this relation, we give a combinatorial interpretation of $Q_n(x+r,y)$ based on Zeng's interpretation of $Q_{n,k}(x)$ in terms of the set $F_{n,k}$ of rooted trees on [n] with root 1 containing k improper edges. We define $F_n^{(r)}$ to be the set of rooted trees on [n] with root 1 for which each child of the root is colored by one of the colors b, w_1, w_2, \ldots, w_r , where b stands for the black color, and w_1, w_2, \ldots, w_r are considered white colors.

We now define a grammatical labeling of a rooted tree $\bar{T} \in F_n^{(r)}$. First, represent an improper edge of \bar{T} by double edges, and denote the resulting tree by \hat{T} . Then the root of \hat{T} is labeled by ax^r , a black vertex is labeled by x and each of the remaining vertices is labeled by w. Moroever, each edge of \hat{T} is labeled by y. In other words, as far as \bar{T} is concerned, a proper edge is labeled by y and an improper edge is labeled by y^2 . For $\bar{T} \in F_n^{(r)}$, we have

$$w(\bar{T}) = ax^{\operatorname{black}(\bar{T}) + r} w^{n - \operatorname{black}(\bar{T})} y^{n + \operatorname{imp}(\bar{T})}, \tag{4.2}$$

where black(\bar{T}) denotes the number of black vertices in \bar{T} .

Using the above labeling scheme, the right-hand side of (4.1) can be expressed as follows.

Theorem 4.2 For $n \ge 0$ and $r \ge 0$,

$$ax^{r}y^{n}w^{n}(r+xw^{-1})Q_{n}(r+xw^{-1},y) = \sum_{\bar{T}\in F_{n+1}^{(r)}}w(\bar{T}).$$
(4.3)

Proof. By (4.2), we see that

$$\begin{split} \sum_{\bar{T} \in F_{n+1}^{(r)}} w(\bar{T}) &= \sum_{\tilde{T} \in F_{n+1}^{(r)}} ax^{\operatorname{black}(\bar{T}) + r} w^{n - \operatorname{black}(\bar{T})} y^{n + \operatorname{imp}(\bar{T})} \\ &= ax^r y^n w^n \sum_{k=0}^{n-1} y^k \sum_{\bar{T} \in F_{n+1}^{(r)}} x^{\operatorname{black}(\bar{T})} w^{-\operatorname{black}(\bar{T})}. \end{split}$$

Given a rooted tree $T \in F_{n,k}$, one can construct a rooted tree \bar{T} in $F_{n,k}^{(r)}$ by assigning the color b to some children of the root 1 and one of the r white colors to each remaining children of the root 1. Thus

$$\begin{split} \sum_{\bar{T} \in F_{n+1,k}^{(r)}} x^{\operatorname{black}(\bar{T})} w^{-\operatorname{black}(\bar{T})} &= \sum_{T \in F_{n+1,k}} \sum_{i=0}^{\deg_T(1)} \binom{\deg_T(1)}{i} (xw^{-1})^i r^{\deg_T(1)-i} \\ &= \sum_{T \in F_{n+1,k}} (r + xw^{-1})^{\deg_T(1)}, \end{split}$$

which can be expressed as $Q_{n,k}(r + xw^{-1})$ according to the interpretation (1.11) of $Q_{n,k}(x)$. It follows that

$$\sum_{\bar{T} \in F_{n+1}^{(r)}} w(T) = ax^r y^n w^n (1 + xw^{-1}) \sum_{k=0}^{n-1} y^k Q_{n,k}(r + xw^{-1}),$$

as claimed.

The following theorem establishes a connection between the grammar H and the sum of weights of rooted trees in $F_n^{(r)}$.

Theorem 4.3 For $n \ge 1$ and $r \ge 0$,

$$D^{n}(ax^{r}) = \sum_{\bar{T} \in F_{n+1}^{(r)}} w(\bar{T}).$$

The proof is similar to that of Theorem 3.1. The operation of adding n as a black child of the root 1 can be described by the substitution rule $a \to axy$ and the operation of adding n as a white child of the root 1 corresponds to the rule $x \to xyw$.

We now give a grammatical derivation of the Berndt-Evans-Wilson-Shor recursion for $Q_{n,k}(x)$, that is, for $n \geq 1$ and $0 \leq k \leq n-1$,

$$Q_{n,k}(1+x) = Q_{n,k}(x) + (n+k-1)Q_{n-1,k}(1+x).$$
(4.4)

Note that $Q_{1,0}(x) = 1$ and $Q_{n,k}(x) = 0$ if $k \ge n$ or k < 0.

Our proof relies on the generating function with respect to the grammar H. For a Laurent polynomial w of the variables in the alphabet V, the exponential generating function of w with respect to D is defined by

$$Gen(w,t) = \sum_{n>0} D^n(w) \frac{t^n}{n!}.$$

We have the following properties:

$$Gen'(w,t) = Gen(D(w),t)$$
(4.5)

$$Gen(w+v,t) = Gen(w,t) + Gen(v,t)$$
(4.6)

$$Gen(wv,t) = Gen(w,t)Gen(v,t), \tag{4.7}$$

where Gen'(w,t) stands for the differentiation of Gen(w,t) with respect to t, and v is also a Laurant polynomial of the variables in the alphabet V, see [2].

We are now in a position to present a grammatical proof of (4.4). It is easily seen that (4.4) follows from the following relation for $n \ge 1$,

$$axy^{n+1}w^{n}(1+xw^{-1})Q_{n}(1+xw^{-1},y)$$

$$= axy^{n+1}w^{n}(1+xw^{-1})Q_{n}(xw^{-1},y)$$

$$+ axw^{n}(1+xw^{-1})\sum_{k=0}^{n-2}(n+k-1)Q_{n-1,k}(1+xw^{-1})y^{n+k+1}.$$
(4.8)

Invoking (4.1) for $n \ge 1$ and r = 0, we obtain that for $n \ge 1$,

$$D^{n}(a) = axy^{n}w^{n-1}Q_{n}(xw^{-1}, y). (4.9)$$

Again, utilizing (4.1) for $n \ge 1$ and r = 1, we find that

$$D^{n}(ax) = axy^{n}w^{n}(1+xw^{-1})Q_{n}(1+xw^{-1},y), (4.10)$$

and so

$$D^{n-1}(ax) = axy^{n-1}w^{n-1}(1+xw^{-1})Q_{n-1}(1+xw^{-1},y).$$
(4.11)

Thus (4.8) can be rewritten as

$$yD^{n}(ax) = yw(1 + xw^{-1})D^{n}(a) + y^{3}w\frac{\partial(D^{n-1}(ax))}{\partial y}.$$
 (4.12)

Expanding (4.11) as

$$D^{n-1}(ax) = axy^{n-1}w^{n-1}(1+xw^{-1})\sum_{k=0}^{n-2}Q_{n-1,k}(1+xw^{-1})y^k,$$

we see that

$$axy \frac{\partial (D^{n-1}(ax))}{\partial a} = xyD^{n-1}(ax),$$

$$xyw \frac{\partial (D^{n-1}(ax))}{\partial x} = ywD^{n-1}(ax),$$

$$yw^2 \frac{\partial (D^{n-1}(ax))}{\partial w} = (n-1)ywD^{n-1}(ax).$$

Notice that

$$D = axy\frac{\partial}{\partial a} + xyw\frac{\partial}{\partial x} + y^3w\frac{\partial}{\partial y} + yw^2\frac{\partial}{\partial w},$$

so that

$$D^{n}(ax) = xyD^{n-1}(ax) + nywD^{n-1}(ax) + y^{3}w\frac{\partial(D^{n-1}(ax))}{\partial y},$$
 (4.13)

and therefore, (4.12) is equivalent to

$$(y-1)D^{n-1}(D(ax)) + (nyw + xy)D^{n-1}(ax) = yw(1+xw^{-1})D^{n-1}(D(a)), (4.14)$$

for $n \geq 1$. In terms of the generating functions, (4.14) can be reformulated as

$$(xy + yw)\operatorname{Gen}(ax, t) + (y - 1 + tyw)\operatorname{Gen}(axyw + ax^{2}y, t)$$

$$= (xy + yw)\operatorname{Gen}(axy, t). \tag{4.15}$$

Let

$$A(t) = (y - 1 + tyw)\operatorname{Gen}(axyw + ax^{2}y, t)$$
$$+ (xy + yw)\operatorname{Gen}(ax, t) - (xy + yw)\operatorname{Gen}(axy, t).$$

Since $D(xw^{-1}) = 0$ as given in (3.8), we have

$$A(t) = (1 + xw^{-1})\operatorname{Gen}(axyw, t) \left(y - 1 + tyw + yw\operatorname{Gen}(y^{-1}w^{-1} - w^{-1}, t) \right).$$

It remains to show that

$$y - 1 + tyw + yw\operatorname{Gen}(y^{-1}w^{-1} - w^{-1}, t) = 0.$$
(4.16)

Observe that

$$D(y^{-1}w^{-1} - w^{-1}) = -y^{-2}w^{-1}y^3w - y^{-1}w^{-2}yw^2 + w^{-2}yw^2 = -1.$$

Hence

$$Gen(y^{-1}w^{-1} - w^{-1}, t) = y^{-1}w^{-1} - w^{-1} - t,$$
(4.17)

which proves (4.16), so that A(t) vanishes. This completes the proof.

5 The Abel Identities

In this section, we present a grammatical approach to the Abel identites. To this end, we establish an expression of $D^n(ax^ry)$ in terms of rooted trees on [n]. Recall that the set of such rooted trees is denoted by R_n .

For a rooted tree $T \in R_n$, we may construct a rooted tree \bar{T} by coloring each child of the vertex 1 by one of the colors b, w_1, w_2, \ldots, w_r . It should be noted that 1 is not

necessarily the root of T. Let $R_n^{(r)}$ denote the set of rooted trees on [n] for which the children of 1 are colored as described above.

We need the following grammatical labeling for a rooted tree $\bar{T} \in R_{n,k}^{(r)}$: First, represent \bar{T} as a rooted tree \hat{T} on $\{0, 1, \ldots, n\}$ with root 0, and represent an improper edge by double edges. Label the vertex 1 by ax^r , label a black vertex by x and label each of the remaining vertices by w. Moreover, each edge in \hat{T} is labeled by y. Thus the weight of \bar{T} is given by

$$w(\bar{T}) = ax^{\operatorname{black}(\bar{T})+r}w^{n-1-\operatorname{black}(\bar{T})}y^{n+\operatorname{imp}(\bar{T})}.$$
(5.1)

Using the same argument as in the proof of Theorem 4.3, we are led to the following relation.

Theorem 5.1 For $n \ge 1$ and $r \ge 0$,

$$D^{n-1}(ax^r y) = \sum_{\bar{T} \in R_n^{(r)}} w(\bar{T}). \tag{5.2}$$

Analogous to Theorem 4.1, there is a connection between $D^n(ax^ry)$ and $Q_n(x,y)$.

Theorem 5.2 For $n \ge 1$ and $r \ge 0$,

$$D^{n-1}(ax^r y) = ax^r y^n w^{n-1} Q_n(r + xw^{-1} - 1, y).$$
(5.3)

In the notation of $Q_n(x, y)$, the relation (2.3) of Dumont and Ramamonjisoa takes the form

$$D^{n-1}(yw) = y^n w^n Q_n(0, y). (5.4)$$

In addition, Dumont [7] obtained grammatical expressions of $Q_n(1, y)$ and $Q_n(-1, y)$: For $n \ge 1$,

$$D^{n}(w) = y^{n}w^{n+1}Q_{n}(1,y), (5.5)$$

$$D^{n}(y) = y^{n+1}w^{n}Q_{n+1}(-1, y). (5.6)$$

It can be checked that by setting a = x = w, the grammar H reduces to the grammar of Dumont and Ramamonjisoa. Meanwhile, (5.4) can be deduced from (5.3) by setting a = x = w and r = 0 and (5.5) can be deduced from (4.1) by setting a = x = w and r = 0.

We remark that (5.6) can also be justified by a grammatical labeling of rooted trees in the set $R_{n,k}[\deg_{\mathbf{T}}(1) = 0]$ of rooted trees in $R_{n,k}$ in which the vertex 1 is a leaf.

For a rooted tree T in $R_{n,k}[\deg_{\mathbf{T}}(1) = 0]$, let \hat{T} denote the tree obtained from T by adding a new root 0 and representing each improper edge by double edges. Label each vertex except for 1 by x and label each edge in \hat{T} by y. Therefore, for a rooted tree in $R_{n,k}[\deg_{\mathbf{T}}(1) = 0]$, we have

$$w(T) = y^{n+k}w^{n-1}. (5.7)$$

Note that the vertex 1 is not endowed with a label. On the other hand, in Shor's construction, it is not allowed to add new vertices as children of the vertex 1. The argument for the proof of Theorem 2.2 implies that for $n \ge 1$,

$$D^{n}(y) = \sum_{T \in R_{n+1,k}[\deg_{T}(1)=0]} w(T).$$

Utilizing the interpretation (1.12) of $Q_{n,k}(x)$, we see that for x = -1,

$$Q_{n,k}(-1) = |R_{n+1,k}[\deg_T(1) = 0]|.$$

Thus it follows from (5.7) that

$$D^{n}(y) = y^{n+1}w^{n} \sum_{k=0}^{n-1} y^{k} Q_{n,k}(-1),$$

which is the right-hand side of (5.6).

The following relations are needed in the grammatical derivations of Abel identities.

Theorem 5.3 For $n \ge 1$,

$$D^{n}(y)|_{y=w=1} = n^{n}, (5.8)$$

$$D^{n}(yw)|_{y=w=1} = (n+1)^{n}, (5.9)$$

$$D^{n}(ax^{r})|_{a=y=w=1} = x^{r}(x+r)(x+r+n)^{n-1},$$
(5.10)

$$D^{n}(ax^{r}y)|_{a=y=w=1} = x^{r}(x+r+n)^{n}.$$
 (5.11)

Proof. In the notation of $Q_n(x,y)$, the relation (1.6) can be rewritten as

$$Q_n(x,1) = (x+n)^{n-1}. (5.12)$$

Setting y = w = 1 in (5.6), we obtain that

$$D^{n}(y)|_{y=w=1} = Q_{n+1}(-1,1),$$

which equals n^n according to (5.12). This proves (5.8). The rest of the relations in the theorem can be obtained from (5.4), (4.1) and (5.3), respectively. This completes the proof.

The classical Abel identity states that for $n \geq 1$,

$$(x+y+n)^n = \sum_{k=0}^n \binom{n}{k} x(x+k)^{k-1} (y+n-k)^{n-k}.$$
 (5.13)

Since y has appeared as a variable in the grammar H, we shall use the following form of (5.13): For $n \geq 1$,

$$(x_1 + x_2 + n)^n = \sum_{k=0}^n \binom{n}{k} x_1 (x_1 + k)^{k-1} (x_2 + n - k)^{n-k}.$$
 (5.14)

Proof of (5.14). Let

$$H': a_1 \to a_1 x_1 y, \ a_2 \to a_2 x_2 y, \ x_1 \to x_1 y w, \ x_2 \to x_2 y w,$$

$$y \to y^3 w, \ w \to y w^2, \tag{5.15}$$

and let D denote the formal derivative associated with the grammar H'. Viewing a_1 as a and x_1 as x and applying (5.10) with respect to H, we get

$$D^{n}(a_{1})|_{a_{1}=y=w=1} = x_{1}(x_{1}+n)^{n-1}.$$
(5.16)

Similarly, invoking (5.11), we obtain that

$$D^{n}(a_{2}y)|_{a_{2}=y=w=1} = (x_{2}+n)^{n}.$$
(5.17)

Moreover, since

$$D(a_1a_2) = a_1a_2(x_1 + x_2)y,$$
 $D(x_1 + x_2) = (x_1 + x_2)yw,$

treating a_1a_2 as a and $x_1 + x_2$ as x, we may apply (5.11) to deduce that

$$D^{n}(a_{1}a_{2}y)|_{a_{1}=a_{2}=y=w=1} = (x_{1} + x_{2} + n)^{n}.$$
(5.18)

Finally, (5.14) is follows from the Leibnitz formula

$$D^{n}(a_{1}a_{2}y) = \sum_{k=0}^{n} \binom{n}{k} D^{k}(a_{1}) D^{n-k}(a_{2}y)$$

along with the relations (5.16), (5.17) and (5.18). This completes the proof.

Riordan [15] obtained a class of generalizations of the classical Abel identity by considering the sum

$$A_n(x, y; p, q) = \sum_{k=0}^{n} \binom{n}{k} (x+k)^{k+p} (y+n-k)^{n-k+q},$$

where $n \ge 1$ and p, q are integers. He found closed forms for the cases when (p, q) lies in

$$\{(-3,0), (-2,0), (-1,0), (0,0), (1,0), (2,0), (-1,-1), (-1,1), (-1,2), (1,1), (1,2), (2,2)\}.$$

These identities for $A_n(x, y; p, q)$ are also called the Abel identities or the Abel-type identities. The original Abel identity corresponds to the case (p, q) = (-1, 0). For (p, q) = (-1, -1), Riordan obtained a closed form for $A_n(x, y; p, q)$, which is stated in the variables x_1 and x_2 : For $n \ge 1$,

$$(x_1 + x_2)(x_1 + x_2 + n)^{n-1} = \sum_{k=0}^{n} \binom{n}{k} x_1 x_2 (x_1 + k)^{k-1} (x_2 + n - k)^{k-1}.$$
 (5.19)

Proof of (5.19). Let H' denote the grammar given by (5.15), and let D denote the formal derivative associated with H'. Using the same reasoning as for the proof of (5.16), we see that

$$D^{n}(a_{2})|_{a_{2}=y=w=1} = x_{2}(x_{2}+n)^{n-1}.$$
(5.20)

Analogous to (5.18), we get the relation

$$D^{n}(a_{1}a_{2})|_{a_{1}=a_{2}=y=w=1} = (x_{1}+x_{2})(x_{1}+x_{2}+n)^{n-1}.$$
 (5.21)

In view of (5.16), (5.20) and (5.21), we are led to (5.19) by applying the Leibnitz formula

$$D^{n}(a_{1}a_{2}) = \sum_{k=0}^{n} \binom{n}{k} D^{k}(a_{1}) D^{n-k}(a_{2}).$$

This completes the proof.

We next consider the case (p,q)=(-2,0) of the Abel identity given by Riordan: For $n\geq 2$,

$$\sum_{k=0}^{n} \binom{n}{k} x_1(x_1+1)(x_1+k)^{k-2}(x_2+n-k)^{n-k}$$

$$= x_1^{-1} \left[(x_1+1)(x_1+x_2+n)^n - nx_1(x_1+x_2+n)^{n-1} \right]. \tag{5.22}$$

Proof of (5.22). Let H' denote the grammar defined as above, and let D denote the formal derivative associated with H'. Set

$$s_1 = a_1 y^{-1} + a_1 x_1^{-1} w,$$

so that $D(s_1) = a_1 x_1$. Analogous to (5.16), we find that for $n \ge 1$,

$$D^{n}(s_{1})|_{a_{1}=y=w=1} = x_{1}(x_{1}+1)(x_{1}+n)^{n-2}.$$
(5.23)

Since

$$s_1 a_2 y = a_1 a_2 + a_1 a_2 x_1^{-1} y w$$

and $x^{-1}w$ is a constant as shown in (3.8), we get

$$D^{n}(s_{1}a_{2}y) = D^{n}(a_{1}a_{2}) + x_{1}^{-1}wD^{n}(a_{1}a_{2}y).$$

By the Leibnitz formula

$$D^{n}(s_{1}a_{2}y) = \sum_{k=0}^{n} {n \choose k} D^{k}(s_{1}) D^{n-k}(a_{2}y),$$

we find that

$$\sum_{k=0}^{n} \binom{n}{k} D^k(s_1) D^{n-k}(a_2 y) = D^n(a_1 a_2) + x^{-1} w D^n(a_1 a_2 y),$$

which yields (5.22) by applying (5.17), (5.18), (5.21) and (5.23). This completes the proof.

We obtain a closed formula for the case (p,q) = (-2,-2).

Theorem 5.4 For $n \ge 1$,

$$\sum_{k=0}^{n} \binom{n}{k} x_1 x_2 (x_1 + 1)(x_2 + 1)(x_1 + k)^{k-2} (x_2 + n - k)^{n-k-2}$$

$$= ((x_1 + x_2)^3 - 3n(x_1 + x_2) - 2n)(x_1 + x_2 + n)^{n-3}$$

$$+ \frac{(x_1 + x_2)^2}{x_1 x_2} (x_1 + x_2 + 1)(x_1 + x_2 + n)^{n-2}.$$
(5.24)

We need the following lemma.

Lemma 5.5 Let D denote the formal derivative associated with the grammar H. For $n \geq 2$,

$$D^{n}(ay^{-1})|_{a=y=w=1} = x(x+1)(x+n)^{n-2} - (x+n)^{n-1}.$$
 (5.25)

For $n \geq 3$,

$$D^{n}(ay^{-2})|_{a=y=w=1} = (x^{3} - 3nx - 2n)(x+n)^{n-3}.$$
 (5.26)

Proof. Since $D(ay^{-1}) = ax - ayw$ and $D(x^{-1}w) = 0$ as given in (3.8), we have

$$D^{n}(ay^{-1}) = D^{n-1}(ax) - x^{-1}wD^{n-1}(axy).$$

Applying (5.10) and (5.11) with r = 1, we get

$$D^{n}(ay^{-1})|_{a=y=w=1} = x(x+1)(x+n)^{n-2} - (x+n)^{n-1}$$

Since

$$D(ay^{-2}) = axy^{-1} - 2aw$$

and

$$D(axy^{-1}) = ax^{2}(1 + x^{-1}w) - axyw,$$

we see that for $n \geq 3$,

$$D^{n}(ay^{-2}) = D^{n-1}(axy^{-1} - 2aw)$$

$$= D^{n-2}(ax^{2}(1 + x^{-1}w) - axyw) - 2x^{-1}wD^{n-1}(ax)$$

$$= (1 + x^{-1}w)D^{n-2}(ax^{2}) - x^{-1}wD^{n-2}(ax^{2}y) - 2x^{-1}wD^{n-1}(ax).$$

In light of (5.10) and (5.11), we find that

$$D^{n}(ay^{-2})|_{a=y=w=1} = x(x+1)(x+2)(x+n)^{n-3} - x(x+n)^{n-2} - 2(x+1)(x+n)^{n-2},$$

which implies (5.26). This complete the proof.

Proof of Theorem 5.4. Assume that H' is the grammar given in (5.15) and D is the formal derivative with respect to H'. Let

$$s_1 = a_1 y^{-1} + a_1 x_1^{-1} w$$

and

$$s_2 = a_2 y^{-1} + a_2 x_2^{-1} w.$$

Clearly, $D(s_1) = a_1 x_1$ and $D(s_2) = a_2 x_2$. It follows from (5.10) that

$$D^{n}(s_{2})|_{a_{2}=y=w=1} = x_{2}(1+x_{2})(x_{2}+n)^{n-2}.$$
(5.27)

By the same argument as that for the proof of (5.18), we deduce from (5.25) and (5.26) that

$$D^{n}(a_{1}a_{2}y^{-1})|_{a_{1}=a_{2}=y=w=1} = (x_{1} + x_{2})(x_{1} + x_{2} + 1)(x_{1} + x_{2} + n)^{n-2} - (x_{1} + x_{2} + n)^{n-1}$$

and

$$D^{n}(a_{1}a_{2}y^{-2})|_{a_{1}=a_{2}=y=w=1} = ((x_{1}+x_{2})^{3} - 3n(x_{1}+x_{2}) - 2n)(x_{1}+x_{2}+n)^{n-3}.$$

Since $D(x_1w^{-1}) = D(x_2w^{-1}) = 0$, we get

$$D^{n}(s_{1}s_{2}) = D^{n}(a_{1}a_{2}(y^{-1} + x_{1}^{-1}w)(y^{-1} + x_{2}^{-1}w))$$

$$= D^{n}(a_{1}a_{2}y^{-2} + a_{1}a_{2}y^{-1}(x_{1}^{-1} + x_{2}^{-1})w + a_{1}a_{2}(x_{1}x_{2})^{-1}w^{2})$$

$$= D^{n}(a_{1}a_{2}y^{-2}) + (x_{1}^{-1} + x_{2}^{-1})wD^{n}(a_{1}a_{2}y^{-1})$$

$$+ (x_{1}x_{2})^{-1}w^{2}D^{n}(a_{1}a_{2}).$$
(5.28)

By the Leibnitz formula

$$D^{n}(s_{1}s_{2}) = \sum_{k=0}^{n} {n \choose k} D^{k}(s_{1}) D^{n-k}(s_{2}),$$

we obtain (5.24) by using (5.23), (5.27) and (5.28). This completes the proof.

We conclude this paper with a one-line grammatical explanation of the Lacasse identity.

Theorem 5.6 For $n \ge 1$,

$$n^{n+1} = \sum_{j=1}^{n} \sum_{k=0}^{n-j} \binom{n}{j} \binom{n-j}{k} j^{j} k^{k} (n-j-k)^{n-j-k}.$$
 (5.29)

Proof. Because of the relation

$$k\binom{n}{k} = n\binom{n-1}{k-1},$$

(5.29) can be rewritten as

$$n^{n} = \sum_{j=1}^{n} \sum_{k=0}^{n-j} {n-1 \choose j-1} {n-j \choose k} j^{j-1} k^{k} (n-k)^{n-k}.$$
 (5.30)

Since $D(y) = y^3 w$, we have

$$D^{n}(y) = D^{n-1}(y^{3}w) = \sum_{i+j+k=n-1} {n-1 \choose i,j,k} D^{i}(y)D^{j}(yw)D^{k}(y).$$
 (5.31)

Invoking (5.8) and (5.9) and setting y = w = 1, we see that (5.31) can be rewritten in the form of (5.30). This completes the proof.

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