

A Telescoping Algorithm for Double Summations

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Abstract

We present an algorithm to prove hypergeometric double summation identities. Given a hypergeometric term $F(n, i, j)$, we aim to find a difference operator $L = a_0(n)N^0 + a_1(n)N^1 + \cdots + a_r(n)N^r$ and rational functions $R_1(n, i, j), R_2(n, i, j)$ such that $LF = \Delta_i(R_1F) + \Delta_j(R_2F)$. Based on simple divisibility considerations, we show that the denominators of R_1 and R_2 must possess certain factors which can be computed from $F(n, i, j)$. Using these factors as estimates, we may find the numerators of R_1 and R_2 by guessing the upper bounds of the degrees and solving systems of linear equations. Our algorithm is valid for the Andrews-Paule identity, the Carlitz's identities, the Apéry-Schmidt-Strehl identity, the Graham-Knuth-Patashnik identity, and the Petkovšek-Wilf-Zeilberger identity.

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1. Introduction

This paper is concerned with double summations of hypergeometric terms $F(n, i, j)$. A function $F(n, k_1, \dots, k_m)$ is called a *hypergeometric term* if the quotients

$$\frac{F(n+1, k_1, \dots, k_m)}{F(n, k_1, \dots, k_m)}, \quad \frac{F(n, k_1+1, \dots, k_m)}{F(n, k_1, \dots, k_m)}, \quad \dots, \quad \frac{F(n, k_1, \dots, k_m+1)}{F(n, k_1, \dots, k_m)}$$

are rational functions of n, k_1, \dots, k_m . Throughout the paper, N denotes the shift operator with respect to the variable n , defined by

$$NF(n, k_1, \dots, k_m) = F(n+1, k_1, \dots, k_m),$$

and Δ_{k_t} denotes the difference operator with respect to the variable k_t , defined by

$$\Delta_{k_t} F(n, k_1, \dots, k_m) = F(n, k_1, \dots, k_t + 1, \dots, k_m) - F(n, k_1, \dots, k_t, \dots, k_m).$$

We also use Δ_i and Δ_j to denote the difference operators with respect to the variables i and j , respectively. For polynomials a and b , we denote by $\gcd(a, b)$ their monic greatest common divisor. When we express a rational function as a quotient p/q , we always assume that p and q are relatively prime unless it is explicitly stated otherwise.

Zeilberger's algorithm [7, 9, 13], also known as the method of *creative telescoping*, is devised for proving hypergeometric identities of the form

$$\sum_k F(n, k) = f(n), \quad (1.1)$$

where $F(n, k)$ is a hypergeometric term and $f(n)$ is a given function. This algorithm has been used to deal with multiple sums by Wilf and Zeilberger [12]. Given a hypergeometric term $F(n, k_1, \dots, k_m)$, the approach of Wilf and Zeilberger is to try to find a linear difference operator L with coefficients being polynomials in n

$$L = a_0(n)N^0 + a_1(n)N^1 + \dots + a_r(n)N^r$$

and rational functions R_1, \dots, R_m of n, k_1, \dots, k_m such that

$$LF = \sum_{l=1}^m \Delta_{k_l}(R_l F). \quad (1.2)$$

The problem of constructing the denominators of R_1, \dots, R_m for the Wilf-Zeilberger approach has not been solved. In an alternative approach, Wegschaider generalized Sister Celine's technique [11] to multiple summations, and proved many double summation identities. In this paper, we provide estimates of the denominators of R_1 and R_2 for double summations. These estimates turn out to be good enough for several double summation identities, including the Andrews-Paule identity which does not seem to be a suitable example for Wegschaider's approach.

To give a sketch of our approach, we first consider Gosper's algorithm for bivariate hypergeometric terms. Suppose that $F(i, j)$ is a hypergeometric term and $p_1/q_1, p_2/q_2$ are rational functions such that

$$F(i, j) = \Delta_i \left(\frac{p_1(i, j)}{q_1(i, j)} F(i, j) \right) + \Delta_j \left(\frac{p_2(i, j)}{q_2(i, j)} F(i, j) \right).$$

We show that under certain hypotheses (Section 2, (H1)–(H3)), the denominators q_1, q_2 can be written in the form

$$\begin{aligned} q_1(i, j) &= v_1(i) v_2(j) v_3(i+j) v_4(i, j) u_1(j) u_2(i, j), \\ q_2(i, j) &= v_1(i) v_2(j) v_3(i+j) v_4(i, j) w_1(j) w_2(i, j), \end{aligned} \tag{1.3}$$

such that v_1, v_2, v_4 and u_2, w_2 are bounded in the sense that they are factors of certain polynomials. Then we apply these estimates to the telescoping algorithm for double summations. Suppose that

$$LF(n, i, j) = \Delta_i (R_1(n, i, j) F(n, i, j)) + \Delta_j (R_2(n, i, j) F(n, i, j)),$$

where

$$R_1(n, i, j) = \frac{1}{d(n, i, j)} \cdot \frac{f_1(n, i, j)}{g_1(n, i, j)}, \quad R_2(n, i, j) = \frac{1}{d(n, i, j)} \cdot \frac{f_2(n, i, j)}{g_2(n, i, j)}$$

and $d(n, i, j)$ is the denominator of $LF(n, i, j)/F(n, i, j)$. We may deduce that g_1, g_2 can be factored in the form of (1.3) such that v_1, v_2, v_4 and u_2, w_2 are bounded. Although we do not have the universal denominators, these bounds can be used to give estimates of the denominators g_1 and g_2 . Then by further guessing the bounds of the degrees of the numerators of R_1 and R_2 , we get the desired difference operator if we are lucky.

Indeed, our approach works quite efficiently for many identities such as the Andrews-Paule identity, Carlitz's identities, the Apéry-Schmidt-Strehl identity, the Graham-Knuth-Patashnik identity, and the Petkovšek-Wilf-Zeilberger identity.

2. Denominators in Bivariate Gosper's Algorithm

For a given bivariate hypergeometric term $F(i, j)$, we give estimates of the denominators of the rational functions $R_1(i, j), R_2(i, j)$ satisfying

$$F(i, j) = \Delta_i (R_1(i, j) F(i, j)) + \Delta_j (R_2(i, j) F(i, j)). \tag{2.1}$$

Let

$$\begin{aligned} R_1(i, j) &= \frac{f_1(i, j)}{g_1(i, j)}, & R_2(i, j) &= \frac{f_2(i, j)}{g_2(i, j)}, \\ \frac{F(i+1, j)}{F(i, j)} &= \frac{r_1(i, j)}{s_1(i, j)}, & \frac{F(i, j+1)}{F(i, j)} &= \frac{r_2(i, j)}{s_2(i, j)}. \end{aligned} \quad (2.2)$$

Dividing $F(i, j)$ on both sides of (2.1) and substituting (2.2) into it, we derive that

$$1 = \frac{r_1(i, j)}{s_1(i, j)} \frac{f_1(i+1, j)}{g_1(i+1, j)} - \frac{f_1(i, j)}{g_1(i, j)} + \frac{r_2(i, j)}{s_2(i, j)} \frac{f_2(i, j+1)}{g_2(i, j+1)} - \frac{f_2(i, j)}{g_2(i, j)}. \quad (2.3)$$

Let

$$u(i, j) = \gcd(s_1(i, j), s_2(i, j)), \quad v(i, j) = \gcd(g_1(i, j), g_2(i, j)),$$

and

$$\begin{aligned} s'_1(i, j) &= s_1(i, j)/u(i, j), & s'_2(i, j) &= s_2(i, j)/u(i, j), \\ g'_1(i, j) &= g_1(i, j)/v(i, j), & g'_2(i, j) &= g_2(i, j)/v(i, j). \end{aligned} \quad (2.4)$$

We concern on those $R_1(i, j), R_2(i, j)$ whose denominators g_1, g_2 satisfy the following three hypotheses.

(H1) For any irreducible factor $p(i, j)$ of $g_1(i, j)$ ($g_2(i, j)$, respectively) and integers h_1, h_2 , $p(i + h_1, j + h_2)$ divides $g_1(i, j)$ ($g_2(i, j)$, respectively) implies $p(i + h_1, j + h_2) = p(i, j)$.

(H2) $\gcd(g'_1(i, j), v(i, j)) = \gcd(g'_2(i, j), v(i, j)) = 1$.

(H3) For any integers h_1, h_2 , $\gcd(g'_1(i + h_1, j + h_2), g'_2(i, j)) = 1$.

Under these hypotheses, we have

Theorem 2.1 *The denominators $g_1(i, j), g_2(i, j)$ can be factored into polynomials:*

$$\begin{aligned} g_1(i, j) &= v_1(i)v_2(j)v_3(i+j)v_4(i, j)u_1(j)u_2(i, j), \\ g_2(i, j) &= v_1(i)v_2(j)v_3(i+j)v_4(i, j)w_1(i)w_2(i, j), \end{aligned}$$

such that

$$\begin{aligned}
v_1(i) &| r_1(i-1, j)s'_2(i-1, j), \\
v_2(j) &| r_2(i, j-1)s'_1(i, j-1), \\
v_4(i, j) &| \gcd(r_1(i-1, j)s'_2(i-1, j), r_2(i, j-1)s'_1(i, j-1)), \\
u_2(i, j) &| \gcd(s_1(i, j)s'_2(i, j), r_1(i-1, j)s'_2(i-1, j)), \\
w_2(i, j) &| \gcd(s_2(i, j)s'_1(i, j), r_2(i, j-1)s'_1(i, j-1)).
\end{aligned}$$

Proof. Substituting (2.4) into (2.3), we get

$$\begin{aligned}
1 &= \frac{r_1(i, j)}{s'_1(i, j)u(i, j)} \frac{f_1(i+1, j)}{g'_1(i+1, j)v(i+1, j)} - \frac{f_1(i, j)}{g'_1(i, j)v(i, j)} \\
&\quad + \frac{r_2(i, j)}{s'_2(i, j)u(i, j)} \frac{f_2(i, j+1)}{g'_2(i, j+1)v(i, j+1)} - \frac{f_2(i, j)}{g'_2(i, j)v(i, j)}.
\end{aligned}$$

That is,

$$\begin{aligned}
&s_1(i, j)s'_2(i, j)g_1(i, j)g'_2(i, j)g_1(i+1, j)g_2(i, j+1) \\
&= f_1(i+1, j)r_1(i, j)s'_2(i, j)g_1(i, j)g'_2(i, j)g_2(i, j+1) \\
&\quad - f_1(i, j)s_1(i, j)s'_2(i, j)g'_2(i, j)g_1(i+1, j)g_2(i, j+1) \\
&\quad + f_2(i, j+1)r_2(i, j)s'_1(i, j)g_1(i, j)g'_2(i, j)g_1(i+1, j) \\
&\quad - f_2(i, j)s_1(i, j)s'_2(i, j)g'_1(i, j)g_1(i+1, j)g_2(i, j+1).
\end{aligned}$$

1. Suppose $p(i, j)$ is an irreducible factor of $v(i, j)$, and for some non-negative integer l , $p^l | v$. Since

$$\gcd(p(i+1, j), f_1(i+1, j)) = \gcd(p(i, j+1), f_2(i, j+1)) = 1,$$

we have

$$p^l(i+1, j) | r_1(i, j)s'_2(i, j)g_1(i, j)g'_2(i, j)g_2(i, j+1)$$

and

$$p^l(i, j+1) | r_2(i, j)s'_1(i, j)g_1(i, j)g'_2(i, j)g_1(i+1, j).$$

There are three cases:

- $p(i, j)$ is a polynomial depending only on i . Then $\gcd(p(i+1, j), g_1(i, j)) = 1$. Otherwise, by hypothesis (H1) we have that $p(i+1, j) = p(i, j)$ is independent of i , which is a contradiction. Similarly, $\gcd(p(i+1, j), g_2(i, j)) = 1$. Since $p(i, j)$ is a polynomial depending only on i , we have

$$\gcd(p(i+1, j), g_2(i, j+1)) = \gcd(p(i+1, j+1), g_2(i, j+1)) = 1.$$

Therefore,

$$p^l(i+1, j) \mid r_1(i, j)s'_2(i, j).$$

- $p(i, j)$ is a polynomial depending only on j . The same discussion leads to

$$p^l(i, j+1) \mid r_2(i, j)s'_1(i, j).$$

- $p(i, j)$ is a polynomial depending both on i and on j . Then either

$$p(i+1, j) = p(i, j+1) \quad \text{or} \quad \gcd(p(i+1, j), p(i, j+1)) = 1.$$

In the former case, $p(i, j)$ is a polynomial of $i+j$ (see [1, Lemma 3] or [8, Lemma 3.3]). In the later case, by hypothesis (H1), we have

$$\gcd(p(i+1, j), g_1(i, j)g'_2(i, j)g_2(i, j+1)) = 1$$

and

$$\gcd(p(i, j+1), g_1(i, j)g'_2(i, j)g_1(i+1, j)) = 1.$$

Thus,

$$p^l(i, j) \mid \gcd(r_1(i-1, j)s'_2(i-1, j), r_2(i, j-1)s'_1(i, j-1)).$$

2. Suppose p is an irreducible factor of g'_1 and $p^l \mid g'_1$ for some non-negative integer l . If $p(i, j) \mid v(i, j+1)$, then $p(i, j-1) \mid v(i, j)$. By hypothesis (H1), $p(i, j-1) = p(i, j)$, which implies $p(i, j) \mid v(i, j)$, contradicting to hypothesis (H2). Noting further that

$$\gcd(f_1(i, j), g_1(i, j)) = \gcd(g'_1(i, j), g'_2(i+h_1, j+h_2)) = 1, \quad \forall h_1, h_2 \in \mathbb{Z},$$

we have

$$p^l(i, j) \mid s_1(i, j)s'_2(i, j)g_1(i+1, j).$$

If $p(i+1, j) \mid v(i, j+1)$, then by hypothesis (H1), $p(i+1, j-1) = p(i, j)$, which implies $p(i, j) \mid v(i, j)$, contradicting to hypothesis (H2). Therefore,

$$p^l(i+1, j) \mid r_1(i, j)s_2'(i, j)g_1(i, j).$$

There are also two cases:

- $p(i, j) = p(i+1, j)$. Then $p(i, j)$ is a polynomial depending only on j .
- $\gcd(p(i, j), p(i+1, j)) = 1$. Then

$$\gcd(p(i, j), g_1(i+1, j)) = \gcd(p(i+1, j), g_1(i, j)) = 1,$$

and hence,

$$p^l(i, j) \mid \gcd(s_1(i, j)s_2'(i, j), r_1(i-1, j)s_2'(i-1, j)).$$

3. Similarly, suppose p is an irreducible factor of g_2' and $p^l \mid g_2'$ for some non-negative integer l . Then either $p(i, j)$ is a polynomial depending only on i or

$$p^l(i, j) \mid \gcd(s_1(i, j)s_2'(i, j), r_2(i, j-1)s_1'(i, j-1)).$$

■

3. Denominators in the Telescoping Algorithm

We are now ready to estimate the denominators of R_1 and R_2 in telescoping algorithm.

As in the case of single summations, the telescoping algorithm for double summations tries to find an operator

$$L = a_0(n) + a_1(n)N + \cdots a_r(n)N^r$$

and rational functions $R_1(n, i, j), R_2(n, i, j)$ such that

$$LF(n, i, j) = \Delta_i(R_1(n, i, j)F(n, i, j)) + \Delta_j(R_2(n, i, j)F(n, i, j)). \quad (3.1)$$

Let

$$\frac{F(n, i+1, j)}{F(n, i, j)} = \frac{r_1(n, i, j)}{s_1(n, i, j)}, \quad \frac{F(n, i, j+1)}{F(n, i, j)} = \frac{r_2(n, i, j)}{s_2(n, i, j)}, \quad (3.2)$$

and $d(n, i, j)$ be the common denominator of

$$\frac{F(n+1, i, j)}{F(n, i, j)}, \quad \dots, \quad \frac{F(n+r, i, j)}{F(n, i, j)}.$$

Then there exists a polynomial $c(n, i, j)$, not necessarily being coprime to d , such that

$$\frac{LF(n, i, j)}{F(n, i, j)} = \sum_{l=0}^r a_l(n) \frac{F(n+l, i, j)}{F(n, i, j)} = \frac{c(n, i, j)}{d(n, i, j)}. \quad (3.3)$$

Note that c is related to the polynomials a_0, a_1, \dots, a_r but d is independent of them.

Now, (3.1) can be written in the form of (2.1):

$$LF(n, i, j) = \Delta_i(R'_1(n, i, j)LF(n, i, j)) + \Delta_j(R'_2(n, i, j)LF(n, i, j)),$$

where

$$R'_1(n, i, j) = R_1(n, i, j) \frac{d(n, i, j)}{c(n, i, j)} \quad \text{and} \quad R'_2(n, i, j) = R_2(n, i, j) \frac{d(n, i, j)}{c(n, i, j)}.$$

This suggests us to assume

$$R_1(n, i, j) = \frac{1}{d(n, i, j)} \frac{f_1(n, i, j)}{g_1(n, i, j)} \quad \text{and} \quad R_2(n, i, j) = \frac{1}{d(n, i, j)} \frac{f_2(n, i, j)}{g_2(n, i, j)}, \quad (3.4)$$

where f_1, g_1 (f_2, g_2 , respectively) are relatively prime polynomials.

Since the following discussion is independent of n , we omit the variable n for convenience. For example, we write $R_1(i, j)$ instead of $R_1(n, i, j)$. Using these notations, we have

Theorem 3.1 *Suppose the polynomials g_1, g_2 in (3.4) satisfy the hypotheses (H1)–(H3). Suppose further that*

$$\gcd(g_1(i, j), d(i+h_1, j+h_2)) = \gcd(g_2(i, j), d(i+h_1, j+h_2)) = 1, \quad \forall h_1, h_2 \in \mathbb{Z}. \quad (3.5)$$

Then $g_1(i, j), g_2(i, j)$ can be factored into polynomials:

$$\begin{aligned} g_1(i, j) &= v_1(i)v_2(j)v_3(i+j)v_4(i, j)u_1(j)u_2(i, j), \\ g_2(i, j) &= v_1(i)v_2(j)v_3(i+j)v_4(i, j)w_1(i)w_2(i, j), \end{aligned}$$

such that

$$\begin{aligned} v_1(i) &\mid r_1(i-1, j)s'_2(i-1, j), \\ v_2(j) &\mid r_2(i, j-1)s'_1(i, j-1), \\ v_4(i, j) &\mid \gcd(r_1(i-1, j)s'_2(i-1, j), r_2(i, j-1)s'_1(i, j-1)), \\ u_2(i, j) &\mid \gcd(s_1(i, j)s'_2(i, j), r_1(i-1, j)s'_2(i-1, j)), \\ w_2(i, j) &\mid \gcd(s_2(i, j)s'_1(i, j), r_2(i, j-1)s'_1(i, j-1)), \end{aligned}$$

where

$$\begin{aligned} s'_1(i, j) &= s_1(i, j) / \gcd(s_1(i, j), s_2(i, j)), \\ s'_2(i, j) &= s_2(i, j) / \gcd(s_1(i, j), s_2(i, j)). \end{aligned} \tag{3.6}$$

Proof. Substituting (3.4) into (3.1) and dividing $F(i, j)$ on both sides, we obtain

$$\begin{aligned} \frac{c(i, j)}{d(i, j)} &= \frac{r_1(i, j)}{s_1(i, j)} \frac{f_1(i+1, j)}{d(i+1, j)g_1(i+1, j)} - \frac{f_1(i, j)}{d(i, j)g_1(i, j)} \\ &\quad + \frac{r_2(i, j)}{s_2(i, j)} \frac{f_2(i, j+1)}{d(i, j+1)g_2(i, j+1)} - \frac{f_2(i, j)}{d(i, j)g_2(i, j)}, \end{aligned} \tag{3.7}$$

i.e.,

$$\begin{aligned} c(i, j) &= \frac{r_1(i, j)d(i, j)}{s_1(i, j)d(i+1, j)} \frac{f_1(i+1, j)}{g_1(i+1, j)} - \frac{f_1(i, j)}{g_1(i, j)} \\ &\quad + \frac{r_2(i, j)d(i, j)}{s_2(i, j)d(i, j+1)} \frac{f_2(i, j+1)}{g_2(i, j+1)} - \frac{f_2(i, j)}{g_2(i, j)}. \end{aligned}$$

Let

$$\begin{aligned} \tilde{r}_1(i, j) &= r_1(i, j)d(i, j), & \tilde{s}_1(i, j) &= s_1(i, j)d(i+1, j), \\ \tilde{r}_2(i, j) &= r_2(i, j)d(i, j), & \tilde{s}_2(i, j) &= s_2(i, j)d(i, j+1). \end{aligned}$$

All discussion in the proof of Theorem 2.1 still holds. Thus, we have

$$\begin{aligned}
v_1(i) &| \tilde{r}_1(i-1, j) \tilde{s}'_2(i-1, j), \\
v_2(j) &| \tilde{r}_2(i, j-1) \tilde{s}'_1(i, j-1), \\
v_4(i, j) &| \gcd(\tilde{r}_1(i-1, j) \tilde{s}'_2(i-1, j), \tilde{r}_2(i, j-1) \tilde{s}'_1(i, j-1)), \\
u_2(i, j) &| \gcd(\tilde{s}_1(i, j) \tilde{s}'_2(i, j), \tilde{r}_1(i-1, j) \tilde{s}'_2(i-1, j)), \\
w_2(i, j) &| \gcd(\tilde{s}_2(i, j) \tilde{s}'_1(i, j), \tilde{r}_2(i, j-1) \tilde{s}'_1(i, j-1)),
\end{aligned} \tag{3.8}$$

where

$$\begin{aligned}
\tilde{s}'_1(i, j) &= \tilde{s}_1(i, j) / \gcd(\tilde{s}_1(i, j), \tilde{s}_2(i, j)), \\
\tilde{s}'_2(i, j) &= \tilde{s}_2(i, j) / \gcd(\tilde{s}_1(i, j), \tilde{s}_2(i, j)).
\end{aligned}$$

Since we have (3.5), we may replace $\tilde{r}_1, \tilde{s}_1, \tilde{r}_2, \tilde{s}_2$ by r_1, s_1, r_2, s_2 in (3.8), respectively. ■

4. A Telescoping Algorithm for Bivariate Hypergeometric Terms

Theorem 3.1 provides us a way to choose the denominators in the telescoping algorithm. Given a hypergeometric term $F(n, i, j)$, we have the following algorithm:

Algorithm EstDen

1. Calculate $r_1, r_2, s_1, s_2, s'_1, s'_2$ defined by (3.2) and (3.6);
2. Set
 - $v_1(i) :=$ the maximal factor of $r_1(i, j) s'_2(i, j)$ depending only on i ;
 - $v_2(j) :=$ the maximal factor of $r_2(i, j) s'_1(i, j)$ depending only on j ;
 - and
 - $v(i) := \gcd(v_1(i-1), v_2(i-1))$;
3. Set
 - $u_1(j) :=$ the maximal factor of $s_1(i, j) s'_2(i, j)$ depending only on j ;
 - $w_1(i) :=$ the maximal factor of $s_1(i, j) s'_2(i, j)$ depending only on i ;
4. Set $u_2(i, j)$ to be the maximal factor of

$$\gcd(s_1(i, j) s'_2(i, j), r_1(i-1, j) s'_2(i-1, j))$$

which depends on i ;

Set $w_2(i, j)$ to be the maximal factor of

$$\gcd(s_1(i, j)s_2'(i, j), r_2(i, j-1)s_1'(i, j-1))$$

which depends on j .

5. Return $g_1(i, j) := v(i)u_1(j)u_2(i, j)$ and $g_2(i, j) := v(i)w_1(i)w_2(i, j)$.

Remark. Let $f(i, j)$ be a polynomial in i, j and a be a new variable. Then the maximal factor of $f(i, j)$ depending only on i can be obtained by

$$\gcd(f(i, j), f(i, j + a)),$$

and the maximal factor of $f(i, j)$ depending on i can be obtained by

$$f(i, j) / \gcd(f(i, j), f(i + a, j)).$$

We are now ready to describe the telescoping algorithm for double summations:

Algorithm BiZeil

1. Using algorithm EstDen to obtain g_1 and g_2 .
2. Set the order r of the linear difference operator L to be zero.
3. For the order r , calculate the common denominator $d(n, i, j)$ of

$$\frac{F(n+1, i, j)}{F(n, i, j)}, \quad \dots, \quad \frac{F(n+r, i, j)}{F(n, i, j)}.$$

(If $r = 0$, then take $d(n, i, j) = 1$.)

4. Set the degrees of f_1 and f_2 to be one more than those of $d \cdot g_1$ and $d \cdot g_2$, respectively.
5. Solve the equation (3.7) by undeterminate coefficients method to obtain a_0, a_1, \dots, a_r and f_1, f_2 .

6. If $a_0 \neq 0$, then return $L, f_1/(d \cdot g_1), f_2/(d \cdot g_2)$ and we are done.
 If $a_0 = 0$ but $\deg f_1 - \deg(d \cdot g_1) \leq 2$, then increase the degrees of f_1 and f_2 by one and repeat Step 5.
 Otherwise, set $r := r + 1$ and repeat the process from Step 3.

Remarks.

1. In most cases, $g_1(i, j)$ and $g_2(i, j)$ can be further reduced by cancelling a factor of degree 1 and 2 from g_1 and g_2 , respectively.
2. In all the following examples except Example 4, the degree of the numerator of R_1 (R_2) is one more than that of the denominator. While in Example 4, the difference is two.

This can be interpreted visually as follows. Let t_1, t_2, t_3, t_4 be the four terms of the right hand side of (3.7) after multiplying the common denominator. In most cases, the leading terms of t_1 and t_2 (t_3 and t_4 , respectively) are cancelled and only these terms are cancelled.

3. There is a trick in Step 5 which accelerates the computation. Given g_1 and g_2 , we may derive part of the factors of f_1 and f_2 by divisibility. For example, suppose (3.7) becomes

$$\frac{c(i, j)}{d(i, j)} = \frac{u_1(i, j)}{v_1(i, j)} f_1(i + 1, j) - \frac{f_1(i, j)}{w_1(i, j)} + \frac{u_2(i, j)}{v_2(i, j)} f_2(i, j + 1) - \frac{f_2(i, j)}{w_2(i, j)},$$

after substituting and simplification. Suppose further that $D(i, j)$ is the common denominator of the above equation. Then we immediately have that $f_1 \cdot D/w_1$ is divisible by $q_1 = \gcd(cD/d, u_1D/v_1, u_2D/v_2, D/w_2)$ and $f_1(i+1, j) \cdot u_1D/v_1$ is divisible by $q_2 = \gcd(cD/d, D/w_1, u_2D/v_2, D/w_2)$, and hence,

$$\frac{q_1}{\gcd(D/w_1, q_1)} \quad \text{and} \quad \frac{q_2}{\gcd(u_1D/v_1, q_2)}$$

are factors of $f_1(i, j)$ and $f_1(i + 1, j)$, respectively.

5. Examples

In the following examples, F denotes the summand of the left hand side of the identity.

Example 1. The Andrews-Paule identity:

$$\sum_{i=0}^n \sum_{j=0}^n \binom{i+j}{i}^2 \binom{4n-2i-2j}{2n-2i} = (2n+1) \binom{2n}{n}^2. \quad (5.1)$$

It was confirmed by Andrews and Paule [2,3] by proving the more general identity

$$\sum_{i=0}^{\lfloor \frac{m}{2} \rfloor} \sum_{j=0}^{\lfloor \frac{n}{2} \rfloor} \binom{i+j}{i}^2 \binom{m+n-2i-2j}{n-2i} = \frac{\lfloor \frac{m+n+1}{2} \rfloor! \lfloor \frac{m+n+2}{2} \rfloor!}{\lfloor \frac{m}{2} \rfloor! \lfloor \frac{m+1}{2} \rfloor! \lfloor \frac{n}{2} \rfloor! \lfloor \frac{n+1}{2} \rfloor!}.$$

Using the algorithm BiZeil, we can deal with (5.1) directly. In fact, we have

$$g_1(i, j) = (2n-2i+1)(n-i+1)(j+1)^2, \quad g_2(i, j) = (2n-2i+1)(n-i+1)(i+1)^2.$$

Cancelling the factors $(n-i+1)$ and $(i+1)^2$ from $g_1(i, j)$ and $g_2(i, j)$, respectively, we obtain

$$\tilde{g}_1(i, j) = (2n-2i+1)(j+1)^2 \quad \text{and} \quad \tilde{g}_2(i, j) = (2n-2i+1)(n-i+1).$$

Finally, we get (in 1 second)

$$(2n+1)F(n, i, j) = \Delta_i R_1 F(n, i, j) + \Delta_j R_2 F(n, i, j),$$

where

$$R_1 = \frac{i^2(6n^2 + 5n + 1 + 6jn^2 + jn - j - in + 2in^2 - 2i - 4j^2n - 2j^2 - 3ij - 4ijn)}{(2n-2i+1)(1+j)^2},$$

$$R_2 = \frac{-2n^2 + 2jn^2 + 6in^2 + 9in + 3jn - 4ijn - 4i^2n - n + j - 3ij + 2i - 4i^2}{(2n-2i+1)},$$

which are the same as given in [11, p. 85]. Summing $i, j = 0, \dots, n$, we get

$$\begin{aligned} & (2n+1) \sum_{i=0}^n \sum_{j=0}^n F(n, i, j) \\ &= \sum_{i=0}^n (R_2 F(n, i, n+1) - R_2 F(n, i, 0)) + \sum_{j=0}^n (R_1 F(n, n+1, j) - R_1 F(n, 0, j)) \\ &= (G(n+1) - G(0)) + R_1 F(n, n+1, n) = (n+1)^2 \binom{2n+1}{n}^2, \end{aligned}$$

where

$$G(i) = \frac{(-2n+i-1)(-4n+2i-1)i}{-1+2i-2n} \binom{4n-2i}{2n-2i}.$$

Example 2. Carlitz's identity [5] (see Also [12, Example 6.1.2]):

$$\sum_i \sum_j \binom{i+j}{i} \binom{n-i}{j} \binom{n-j}{n-i-j} = \sum_{l=0}^n \binom{2l}{l}.$$

We have

$$g_1(i, j) = (j+1)^2(-n+j), \quad g_2(i, j) = (i+1)^2(-n+i).$$

Cancelling the factors $(-n+j)$ and $(i+1)(-n+i)$, we obtain

$$\tilde{g}_1(i, j) = (j+1)^2 \quad \text{and} \quad \tilde{g}_2(i, j) = i+1.$$

Noting that for recurrence of order 2, $d(i, j) = (-n+i-1+j)^2(-n+i-2+j)^2$, we finally get (in 2 seconds)

$$LF(n, i, j) = \Delta_i R_1 F(n, i, j) + \Delta_j R_2 F(n, i, j),$$

where

$$L = (4n+6) - (8+5n)N + (n+2)N^2,$$

and

$$R_1 = \left(-i^2(-n+i-1)(36-10ji^2n-13j^2ni+60j^2+60ji-2i^2-38j^2i-8ji^2+10i^3+36n^3-11in^3-14jn^3-2i^4-92jn^2+8i^2n-80in+5j^2n^2+8j^2i^2+88jin+42j^2n-172jn+24jin^2+5i^2n^2+3i^3n-54in^2+88n^2+4j^3n-90j+6j^3-40i+5n^4+90n) \right) / ((-n+i-1+j)^2(-n+i-2+j)^2(j+1)^2),$$

$$R_2 = \left((64-19ji^2n-6j^2ni+14j^2+74ji+54i^2-10j^2i-36ji^2+2i^3+39n^3-16in^3-9jn^3-4i^4+6ji^3-53jn^2+50i^2n-176in+4j^2n^2+4j^2i^2+5n^4+83jin+16j^2n-100jn+22jin^2+11i^2n^2+4i^3n-93in^2+112n^2-60j-108i+140n)(-n-1+j) \right) / ((-n+i-2+j)^2(-n+i-1+j)^2).$$

Example 3. Carlitz's identity [4] (see Also [12, Example 6.1.3]):

$$\begin{aligned} \sum_i \sum_j \binom{i+j}{i} \binom{m-i+j}{j} \binom{n-j+i}{i} \binom{m+n-i-j}{m-i} \\ = \frac{(m+n+1)!}{m!n!} \sum_k \frac{1}{2k+1} \binom{m}{k} \binom{n}{k}. \end{aligned}$$

We have

$$g_1(i, j) = (n-j+i)(1+j)^2, \quad g_2(i, j) = (m-i+j)(i+1)^2.$$

Cancelling the factors $(1+j)$ and $(i+1)^2$, we obtain

$$\tilde{g}_1(i, j) = (n-j+i)(1+j) \quad \text{and} \quad \tilde{g}_2(i, j) = m-i+j.$$

Noting that for recurrence of order 2, $d(i, j) = (-n+j-1)^2(-n+j-2)^2$, we finally get (in 37 seconds)

$$LF(n, i, j) = \Delta_i R_1 F(n, i, j) + \Delta_j R_2 F(n, i, j),$$

where

$$\begin{aligned} L = 2(m+3+n)(2+m+n)^2 - \\ (3m+2nm+4n^2+14+15n)(n+m+3)N + (2n+5)(n+2)^2 N^2, \end{aligned}$$

and the denominators of R_1, R_2 are $d(i, j)\tilde{g}_1(i, j)$ and $d(i, j)\tilde{g}_2(i, j)$, respectively. The degrees of denominators and numerators of R_1, R_2 are both less than those given in [12].

Example 4. The Apéry-Schmidt-Strehl identity [10]:

$$\sum_i \sum_j \binom{n}{j} \binom{n+j}{j} \binom{j}{i}^3 = \sum_k \binom{n}{k}^2 \binom{n+k}{k}^2.$$

We have

$$g_1(i, j) = (-j-1+i)^3 \quad \text{and} \quad g_2(i, j) = (i+1)^3.$$

Cancelling the factors $(-j-1+i)$ and $(i+1)^2$, we obtain

$$\tilde{g}_1(i, j) = (-j-1+i)^2 \quad \text{and} \quad \tilde{g}_2(i, j) = i+1.$$

Noting that for recurrence of order 2, $d(i, j) = (n + 2 - j)(n + 1 - j)$, we finally get (in 1 second)

$$LF(n, i, j) = \Delta_i R_1 F(n, i, j) + \Delta_j R_2 F(n, i, j),$$

where

$$L = (n + 1)^3 - (3 + 2n)(17n^2 + 51n + 39)N + (n + 2)^3 N^2,$$

and

$$R_1 = (-2i^2(3+2n)(-10+30j^2-49n^2-j^3-4n^4-24n^3-2n^2i^2+n^2i-6ni^2+3ni+3nji+n^2ji+3j^2i^2-3j^3i+3ji-4i^2-2j^2i-2ji^2+11n^2j^2+6n^2j+33nj^2+18nj-6j^4+2i+15j-39n)) / ((n+2-j)(n+1-j)(-j-1+i)^2),$$

$$R_2 = (2(-j+i)(3+2n)(-8n^2i-4n^2i^2-4n^2ji+4n^2j+4n^2j^2+12nj-12nji-24ni+12nj^2-12ni^2+12j^2-4ji^2+j^3+6j^2i^2-3j^4+8j+5j^2i-8i^2+3j^3i-16i-16ji)) / ((n+2-j)(n+1-j)(i+1)).$$

The rational functions R_1, R_2 are simpler than those given in [10]. The operator L was used by Apéry in his proof of the irrationality of $\zeta(3)$ and Chyzak and Salvy obtained it using Ore algebras [6].

Example 5. The Strehl identity [10]:

$$\sum_i \sum_j \binom{n}{j} \binom{n+j}{j} \binom{j}{i}^2 \binom{2i}{i}^2 \binom{2i}{j-i} = \sum_k \binom{n}{k}^3 \binom{n+k}{k}^3. \quad (5.2)$$

We have

$$g_1(i, j) = (j+1-i)^3, g_2(i, j) = (-3i-3+j)(-3i-2+j)(-3i-1+j)(i+1)^3.$$

Cancelling the factor $(-3i-3+j)(-3i-2+j)$ from g_2 , we obtain

$$\tilde{g}_1(i, j) = (j+1-i)^3 \quad \text{and} \quad \tilde{g}_2(i, j) = (-3i-1+j)(i+1)^3.$$

Noting that for recurrence of order 6,

$$d(i, j) = (n+1-j)(n+2-j)(n+3-j)(n+4-j)(n+5-j)(n+6-j),$$

we finally get (in 2510 seconds)

$$LF(n, i, j) = \Delta_i R_1 F(n, i, j) + \Delta_j R_2 F(n, i, j),$$

where L is a linear difference operator of order $r = 6$ and the denominators of R_1, R_2 are $d(i, j)\tilde{g}_1(i, j)$ and $d(i, j)\tilde{g}_2(i, j)$, respectively. The operator L is the same as the operator obtained by applying Zeilberger's algorithm to the right hand side of (5.2).

Example 6. The Graham-Knuth-Patashnik identity [7, p. 172]:

$$\sum_j \sum_k (-1)^{j+k} \binom{j+k}{k+l} \binom{r}{j} \binom{n}{k} \binom{s+n-j-k}{m-j} = (-1)^l \binom{n+r}{n+l} \binom{s-r}{m-n-l}. \quad (5.3)$$

We have

$$g_1(j, k) = (k+1)(k+l+1), \quad g_2(j, k) = (j+1)(j+1-l).$$

Cancelling the factor $(j+1)(j+1-l)$ from g_2 , we obtain

$$\tilde{g}_1(j, k) = (k+1)(k+l+1) \quad \text{and} \quad \tilde{g}_2(j, k) = 1.$$

Noting that for recurrence of order 1 (with respect to r), $d(j, k) = r - j + 1$, we finally get (in 8 seconds)

$$LF(r, j, k) = \Delta_j R_1 F(r, j, k) + \Delta_k R_2 F(r, j, k),$$

where

$$L = (r+n+1)(n+s+l-m-r) + (r-l+1)(r-s)R$$

is a linear difference operator with respect to the variable r and the denominators of R_1, R_2 are $d(j, k)\tilde{g}_1(j, k)$ and $d(j, k)\tilde{g}_2(j, k)$, respectively. Then (5.3) follows from the evaluation of the initial value ($r = 0$) by Zeilberger's algorithm:

$$\sum_k (-1)^k \binom{k}{k+l} \binom{n}{k} \binom{s+n-k}{m} = (-1)^l \binom{n}{n+l} \binom{s}{m-n-l}.$$

Example 7. The Petkovšek-Wilf-Zeilberger identity [9, p. 33]:

$$\sum_r \sum_s (-1)^{n+r+s} \binom{n}{r} \binom{n}{s} \binom{n+s}{s} \binom{n+r}{r} \binom{2n-r-s}{n} = \sum_k \binom{n}{k}^4. \quad (5.4)$$

We have

$$g_1(r, s) = (n+r)(n+1-r)(s+1)^2, \quad g_2(r, s) = (n+r)(n+1-r)(r+1)^2.$$

Cancelling the factors $s+1$ and $(r+1)^2$, we obtain

$$\tilde{g}_1(r, s) = (n+r)(n+1-r)(s+1) \quad \text{and} \quad \tilde{g}_2(r, s) = (n+r)(n+1-r).$$

Noting that for recurrence of order 2, $d(r, s)$ equals

$$(n+1)(n+2)(n+1-r)(n+2-r)(n+1-s)(n+2-s)(n-r-s+1)(n+2-r-s),$$

we finally get (in 35 seconds)

$$LF(n, r, s) = \Delta_r R_1 F(n, r, s) + \Delta_s R_2 F(n, r, s),$$

where

$$L = 4(4n+5)(4n+3)(n+1) + 2(2n+3)(3n^2+9n+7)N - (n+2)^3 N^2$$

is a linear difference operator and the denominators of R_1, R_2 are $d(r, s)\tilde{g}_1(r, s)$ and $d(r, s)\tilde{g}_2(r, s)$, respectively. The recursion is the same as that obtained by applying Zeilberger's algorithm to the right hand side of (5.4).

Example 8. Calculate

$$f(n) = \sum_{i=0}^n \sum_{j=0}^i (i^2 + j^3) \binom{n}{i} \binom{i}{j}.$$

We have

$$g_1(i, j) = (j+1)(i^2 + j^3), \quad g_2(i, j) = (i+1-j)(i^2 + j^3).$$

Cancelling the factors $j+1$ and $i^2 + j^3$, we obtain

$$\tilde{g}_1(i, j) = i^2 + j^3 \quad \text{and} \quad \tilde{g}_2(i, j) = i+1-j.$$

Noting that for recurrence of order 1, $d(i, j) = n+1-i$, we finally get (in 1 second)

$$LF(n, i, j) = \Delta_i R_1 F(n, i, j) + \Delta_j R_2 F(n, i, j),$$

where

$$L = 3(n^2 + 20n + 27)(n+1) - (n^2 + 18n + 8)nN$$

and the denominators of R_1, R_2 are $d(i, j)\tilde{g}_1(i, j)$ and $d(i, j)\tilde{g}_2(i, j)$, respectively. Solving $Lf(n) = 0$, we immediately get

$$f(n) = 3^{n-3}n(n^2 + 18n + 8).$$

Using the package `MultiSum.m` by K. Wegschaider and A. Riese, we only find recurrences of order greater than 2, which are not easy to solve.

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